(19) Europäisches Patentamt
European Patent Office
Office européen des brevets

(11) EP 0 817 422 A2

(12) EUROPEAN PATENT APPLICATION

(43) Date of publication:
07.01.1998 Bulletin 1998/02

(51) Int. Cl.[6]: **H04L 12/24**

(21) Application number: **97117274.7**

(22) Date of filing: **18.08.1993**

(84) Designated Contracting States:
BE CH DE DK ES FR GB GR IE IT LI NL

(30) Priority: **28.08.1992 SE 9202488**
05.02.1993 SE 9300363

(62) Document number(s) of the earlier application(s) in accordance with Art. 76 EPC:
93919758.8 / 0 627 143

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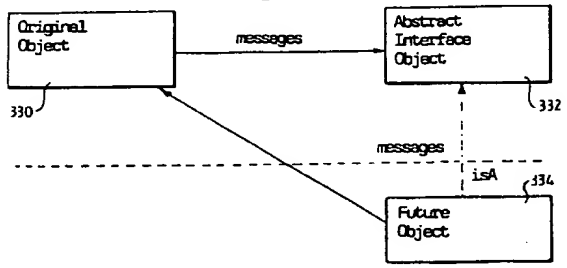
Remarks:
This application was filed on 06 - 10 - 1997 as a divisional application to the application mentioned under INID code 62.

(54) A method for implementing a managed object in a network management system

(57) The invention relates to a method for implementing a managed object in a subsystem of a managed system in a management network with at least one managing system and at least one managed system, for telecom or open systems, said managed system consisting of subsystems forming each a part of the managed system including one or more managed objects.

The managed objects are implemented in the subsystem, uncoordinatedly with respect to other subsystems, in such a way that they can be connected to and transmit messages to other objects in other subsystems, and without knowing the type of objects in the other subsystems. This is obtained by designing a first object (330) for co-operation with an abstract object (332) defining an interface consisting of unimplemented methods, which may be called by the first object, and letting, at later design of a second object (334) unknown to said first object and intended to co-operate with the first object, the other object inherit the abstract object and implement the inherited methods, so that the first object at co-operation with the second object will consider it as being of said abstract type.

Fig.7



Description

Technical field

The present invention relates to a method for implementing a managed object in a management network with at least one managing system and at least one managed system, for telecom or open systems, said managed system consisting of subsystems forming each a part of the managed system including one or more managed objects.

With a "management network with at least one managing system and at least one managed system" is meant that the management network can include at least one managing system which can manage one or more managed systems, which likewise can form part of the management network.

With an open system is meant a system of the kind, which is defined in Reference Model of Open Systems Interconnection (OSI) for CCITT Applications, Rec. X.200.

To perform management activities in a management domain there must be at least one manager which is responsible for the management of the resources. A resource is something which includes concepts and ability of the domain. An example of a domain is a telecom network, where the resources are switches, trunks etc. and management units are e.g. operator tools and managing systems for the network.

For operation of telephone networks each individual company has used a number of different systems for operation and maintenance. CCITT has developed a standard model for operation and maintenance in telephone networks, called TMN (Telecommunication Management Networks). The basic principle of TMN is to indicate an organized network structure, which admits connection of various managing systems to telecom equipment. This is achieved by use of standardized protocols and interfaces. The telephone companies and other operators will require that future telecom networks are adapted to TMN.

CCITT has a recommendation for this under development, the M.3000-serie.

TMN considers all network nodes as network elements (NE). These network elements are made as telephone switches and transmission or transport network products.

The functional structure of TMN comprises

- management functions (OSF, Operations Support Functions), which manage application programs available for users, such as management functions for "Business Management" and service and network administration;
- data communication functions (DCF; Data Communications Functions), which manage data communication between the managing systems OSS and the managed systems NE;
- mediation functions (MF, Mediation Functions), which convert information (between e.g. managed objects), manage data, concentrate, reduce and edit, make decisions relating to e.g. threshold limits and store data, which identify equipment and networks;
- network element functions (NEF, Network Element Functions), which manage telecom processes as switching functions and transmission, and take part in management processes for telecommunication, as fault localisation and protection connections;
- functions for interface adaption (QAF, Q-Adapter Functions), which perform conversion of interfaces from non standard to standard;
- work station functions (WSF, Work Station Functions), which constitute the user terminals of TMN, show information and assist management technicians to manage the network.

TMN also includes an interface, called Q3. Q3 is besides being a communication protocol also an information model, which comprises data schemes, operations and notifications. The exact details of the Q3 interface and its protocols appear from the CCITT recommendations Q961 and Q962.

TMN considers all physical and logical objects as managed objects, which in TMN are referred to as MO (Managed Objects), a denomination which will be used alternately also here henceforth. Managed objects are data images of such physical or logical resources as wires, circuits, signal terminals, transmission routes, events logs, alarm reports etc..

A specific relationship occurs between a resource and a managed object. A resource can be related to one or more MO or none at all. On the other hand an MO can be related to one or more resources or none at all. This relationship is important if an MO is affected by some form of operation or maintenance activity. An MO must not be removed before the functions to which it is subordinated have themselves been removed. This information model is based upon object orientation and the relation concept.

The managing system (OSS- Operation Support system) treats network elements and subordinated managing systems as a collection of managed objects in an imagined data base MIB (Management Information Base). These managed objects are constituted by instances of an MO class, such as a number of signal terminals of the same type. Each terminal will thus constitute an instance of the class signal terminals.

In TMN there also exists the concept MIM (Management Information Model), which refers collectively to all information related to managed objects. MIM is a model of all attributes, relations, operations and notifications referable to managed objects. To be able to search for MO-instances a management information tree MIT (Management Information Tree) is used. This tree structure starts in the network and indicates network ele-

ments and subscribers or equipment.

For operation and maintenance each separate unit or resource, about which information is required, or the operation of which is influenced from the outside of the system, is represented by a managed object. Each information exchange referable to the management of operations or units, which must be influenced or which must report something (such as set up of data, assignment of name, or management of alarms) is done in the form of operations on, or notes from managed objects.

A managed object includes attributes and can also have a relation to other managed objects. A number of different operations can be directed towards a managed object and events can be generated by such objects.

Below an account of deficiencies of TMN will be given.

Network elements and subordinated managing systems are managed by a managing system by means of operations towards the managed objects in the managed systems and supervision of notifications transmitted by the managed system.

Allowed operations towards managed objects in the managed system are determined by an information model of the managed system, together with the notifications, which can be transmitted. The information model states:

- which classes of managed objects that are defined,
- which operations the managed objects in these classes accept,
- which notifications that the managed objects are expected to transmit,
- how many instances that can be created or removed, in each class of managed objects,
- dependences between managed objects, e.g. that one managed object requires existence of another one,
- dependences in a managed object, e.g. that a specific attribute value of one attribute is only allowed if another attribute is set to a specific value or if the managed object only can be removed if it is in a specific state,
- the purpose and intention with managed objects and their notifications.

To be meaningful the managing system must know the information model of the managed system. This is in TMN called "shared management knowledge".

At a change of the information model the management model must be updated in accordance with this change. In conventional systems these changes are made by:

- Definition of the new information model. This is made by specifications for managed objects written as GDMO/ASN.1 templates (GDMO - Guidelines for the Definition of Managed Objects according to CCITT rec. X.722 ISO/IEC 10165-4) and ER-diagrams (Entity-Relationship diagrams) for the managed objects. The specifications for the managed objects state formally (machine readable) syntax (e.g. operations and notifications) for the managed object.

All other parts of the information model, as dependences, the number of instances etc., are stated informally as comments in natural language.

- Implementation and verification of the new information model in the managing system and the managed system.
- Confirmation that the managing system and the managed systems are adapted to the same information model by performing accepted test sequences.
- Updating of the network consisting of the managing system and the managed system with this new version of the information model.

That just mentioned above results in a number of problems:

Firstly, the development of managing systems and managed systems must be coordinated, which leads to higher development costs and/or a delayed introduction of new services on the market.
Secondly, the absence of formalism with respect to the specifications of the managed systems, makes implementation, verification and acceptance of both managing system and managed systems to a difficult and time demanding task, since the interpretation of the specifications is open to discussion.
Thirdly, updating of networks must be planned and carried out carefully, as there are dependences between different versions of the managing systems and managed systems. This involves a delayed introduction of new services in the network.

The purpose of management according to the TMN model is to be a framework for standardization of management for telecom networks and open systems. The management structure strongly influences the management paradigm for new system architectures. There are strong reasons to assume the management paradigm according to the TMN model for the whole management and not only for domains subjected to standardization. The main reason for this is that it is desirable to be able to develop and design management functions in a uniform way.

Summary of the invention

It is a purpose of the invention to design a managed object, which should be able to communicate with unknown future objects. It should be possible for the original object to be related to and to co-operate with an

unknown object without modification, or even reloading of the original object.

Said purpose is attained by the method defined by way of introduction being characterized in that the managed objects are implemented in the subsystem, uncoordinatedly with respect to other subsystems, in a way that they can be connected to and transmit messages to other objects in other subsystems, and without knowing the type of objects in the other subsystems, by designing a first object for co-operation with an abstract object defining an interface consisting of unimplemented methods, which may be called by the first object, and having, at later design of a second object unknown to said first object and intended to co-operate with the first object, the second object to inherit the abstract object and implement the inherited methods, so that the first object at co-operation with the second object will consider it as being of said abstract type.

Advantageous embodiments of the invention have obtained the characterizing features stated in the respective dependent claims.

Description of the drawings.

A number of embodiments of the invention will now be described below in greater detail with reference to the enclosed drawings, on which

Figures 1-9 in block diagrams illustrate problems, which can appear at the use of technology according to the state of the art in connection with the design of managing systems, where
Figures 1-4 concern problems, which can appear in connection with reuse of library components in the managed objects in managed systems,
Figures 5 and 6 refer to problems, which can appear at the implementation of a managed system in a layered system architecture,
Figures 7-10 illustrate how the problems according to Figures 1-6 can be generally solved by design of a specific type of object, which should be able to co-operate with unknown future managed objects,
Figures 11-14 specify design of objects according to Figures 10-14, the object type being assumed to belong to a platform system,
Figures 15-20 show the implementation in the program code C++ of dependences between the designed objects according to Figures 11-14.

Description of embodiments

A managed system consists of several subsystems, more exactly a system platform and a number of applications. According to the invention these subsystems are developed uncoordinatedly and are loaded separately.

At the development of an application - or a system platform - there are libraries with reusable components. These components should be incorporated and combined in different ways in the different subsystems.

The two different problems described above have certain properties in common. There are system components, which are developed uncoordinatedly, but still there are dependences between the components. To keep such dependences it is necessary that the components should be able to communicate with other components which are separately developed and even loaded.

The design process described below in greater detail is usable in two different contexts with similar problems.

The first problem refers to reuse of library components.

At design of a managed system there are libraries with reusable components, which can be incorporated in the objects of the managed system. As is illustrated in Figure 1 there are first a main library 260 including components with basic functionality, such as MOstateComponent 262 and WorkingStateComponent 264. These components include state attributes common to many different types of managed objects. The functionality of these basic components is reused by other components with a more specialized functionality. In the traffic management library 266 there is e.g. a StatePropagationComponent 268, which has the ability to transfer state transitions in MOstateComponent 262 to the related dependent object 270.

It should be emphasized that components as well as the managed objects are supposed to be implemented in an object oriented language. Both the components and the managed objects are thus really objects. But the components can not exist as identifiable objects by themselves in an application. They can only form a part of managed objects.

A straightforward way to implement components, which depend on other basic components, e.g. the components in the fault handling and the traffic management libraries in Figure 1, is to include the basic components either by heritage or aggregating. This causes however many problems. This depends upon the fact that several components, e.g. FaultHandlingComponent 272 and ResourceManagementComponent 274, which inherit or aggregate the same basic component, e.g. MOstateComponent 262, can be included in the same managed object-Route 276, compare Figure 2. This implies that there will be several copies of the basic component in the managed object, which will cause consistency problems if the basic component includes data, which is visible external to the components. Data of MOstateComponent should in other words be available for another functionality in the managed object than those components in which MOstateComponent is included.

There is a way to avoid the problem described above.

MOstateComponent should not be inherited or aggregated in components which depend on this com-

ponent. Instead the dependent components should include reference attributes to MOstateComponent. Thus the latter will be a component of its own in those objects in which it is included. Components which need to access MOstateComponent will then each include a reference to the same instance of the MO-state component, compare Figure 3.

As has been mentioned earlier a component, such as ResourceManagementComponent 274 and FaultHandlingComponent 272, can include a functionality which is triggered by state changes in MOstateComponent 262. The latter must therefore be able to transmit state change messages to other components. One problem is however that the receivers of these messages are unknown at the design of MOstateComponent. MOstateComponent 262 must in some way be able to communicate, cf. arrows 280, 282, with an arbitrary later appearing component, compare Figure 4.

The second problem refers to a layered system architecture.

A managed system can be implemented in a layered architecture. More exactly, there are system platforms with basic functions, which can be reused by different applications, compare Figure 5. The system platform 290 i.a. includes different kinds of common telecom services. The application can thus delegate many tasks to the system platform. The system platform 290 and the applications 292 are loaded separately. It must be possible to load an application and link it to the system platform without reloading the platform.

Both the system platform 290 and the applications 292 include managed objects. The objects 294 and 296 in the platform 290 can e.g. represent resources, as switches, transmission routes, trunks etc. The objects 298,300 and 302,304 in the applications 292.1 resp 292.2 can be related to and co-operate with the resources in the system platform, compare Figure 6. An object in an application can delegate some tasks to an object in the system platform. Objects in the applications may thus depend on related objects in the system platform to be able to work. As the system platform is known at the development of an application it is no problem to design the objects in the applications for communicating with the objects in the system platform. As has been mentioned earlier the objects in the applications can however be dependent on the objects in the system platform. This implies that objects in the system platform should be able to transfer state changes, as has been described earlier above, to objects in the applications. This implies that objects in the platform should be able to be related to and take initiative to calling objects, which are unknown at the design of the platform system.

Above two cases with similar problems have been described. Here will now follow a description of how these problems can be solved.

More exactly, the real problem in each of the two cases is to design an object, which should be able to communicate with unknown future objects. It must be possible for the original object to be related to and to co-operate with an unknown object without modification, or even reloading of the original object.

This problem can be solved with the design principle illustrated in Figure 7. The original object 330 is designed to co-operate with an abstract object 332. The abstract object defines an interface consisting of unimplemented methods. These are the methods, which can be called by the original object. At the design of an object 334 designated to co-operate with the original object it must inherit the abstract object 332 and implement its inherited methods. At the co-operation with an unknown type of object the original object 330 will consider this as being said abstract type 332. At call of a method defined in the interface of the abstract object 332 the call will be delegated to implementation in the real object 330 via late binding.

Avoidance of reloading of the original object at loading of the future objects is achieved by dynamic linking.

Figure 8 illustrates how the just described design can be used for the design of basic library components. MOstateComponent 340 is designed to communicate with objects, which inherit an abstract object: MOstateSubscriber 342. This abstract object defines methods called by MOstateComponent when a state change occurs. If a component should subscribe to state changes in MOstateComponent, it must inherit MOstateSubscriber 342 and implement the answer to the respective state transition notifications. The subscribing component must of course also state itself as a subscriber to MOstateComponent.

The principles of design described with reference to Figure 7 can also be used for designing objects in a system platform to co-operate with application specific objects. This is illustrated in Figure 9, where object PlatformObject1 350 is designed to co-operate with object 352 in application system 354. An object in an application, which should co-operate with the object Platformobject1 350 must inherit the abstract object Interfaceobject 356. Further there is a data base relationship between PlatformObject1 350 and InterfaceObject 356. This implies that the objects which inherit InterfaceObject 356 can set up relations with PlatformObject1 350. To make it possible to load applications without reloading of the system platform dynamic linking is used.

Often an object in an application system is state dependent of another object in the system platform. This design makes it possible to implement such state dependences in essentially the same way as for objects which are known to each other when they are implemented, as will be described in greater detail further below.

Dependences between uncoordinated objects can be specified and implemented in essentially the same way as between coordinated objects, which has been described earlier above. Here the same examples will

be used to show the design and implementation of dependences between uncoordinated objects. It is however assumed here that the object type ResourceB belongs to a platform system. Thus, with reference to Figure 10, various kinds of objects in various application systems can be related to and co-operate with the object ResourceB 370. The object type ResourceB 370 is related to UserObject 372 via a data base relationship. ResourceA 374 inherits UserObject 372 to be able to be related to ResourceB 370. In the following it will be shown how these objects are specified and implemented. The same pseudo syntax as above will be used.

The specification of the object ResourceB is found in Figure 11. The specification includes two state attributes: admState and opState, a method to allocate instances of ResourceB and a precondition which states that to enable erase of a ResourceB object it must be locked. This is essentially the same as earlier. The only difference is that in this example ResourceB is related to UserObject instead of ResourceA. The relation is established via a reference attribute userRef, line 6.

In Figure 12 the object UserObject is specified. It includes the state attribute admState. It can however be noted that this attribute is specified as being purely virtual, line 19. This implies that the UserObject in fact will not include the attribute admState. The interface of UserObject will however include unimplemented write and read methods for this attribute. The real implementation of these virtual methods will appear in the respective subtypes of UserObject. Besides there are another attribute resourceBstate, which is assumed to keep the value of opState in the related object ResourceB. Further, the UserObject has a reference attribute Bref, which is used to establish relations with objects of the type resourceB, line 21.

A dependence scheme, which specifies end conditions, including both the object UserObject and the object ResourceB is found in Figure 13. There is an end condition, lines 29-31, which states that admState of UserObject depends on admState in ResourceB in such a way that an object UserObject can not be locked up if the related object ResourceB is locked.

In Figure 13 it is further specified that when a value of opState in an object ResourceB is changed, the new value should be propagated to the related instance of UserObject. In Figure 13 there is no end condition specified which relates to opState, it is only stated that changes of opState in ResourceB should be propagated to the attribute resourceBstate in UserObject, lines 36 and 37. This implies that all subtypes of UserObject are not necessarily opState dependent on ResourceB, but each subtype of UserObject can manage the opState dependence of ResourceB in its own way.

In Figure 14 the specification of ResourceA is shown. ResourceA is specified as being a subtype of UserObject (line 41). This implies that ResourceA inherits the attributes, methods, conditions and the propagations in UserObject. All the purely virtual specifications in UserObject must be specified and implemented in ResourceA. As appears from Figure 7 the value of opState in ResourceA is derived from the attributes internalOpState and resourceBstate (which is inherited from UserObject), lines 46-49.

As has been mentioned earlier the implementation is essentially the same as when the objects belong to the same system. The difference is that a certain part of the functionality of the dependent object which uses ResourceB is implemented in UserObject. As earlier the code can automatically be generated from the specifications by a compiler. The declaration file which is generated from the specification of UserObject in Figure 12 is found in Figure 15.

The method checkConsistency in UserObject checks the dependence between the admState attributes, which are stated in the dependence scheme in Figure 13. The implementation of this method is shown in Figure 16. To check this condition the value of admState in UserObject must be read. This explains why there must be a purely virtual specification of admState in UserObject, even if this attribute is not really included in UserObject.

The declaration file which is generated by the specification of ResourceB in Figure 11 is found in Figure 17. The implementation of the method propagateOpstate is however somewhat different, compare Figure 8. The new value is propagated to the attribute ResourceBstate in the related userObject instance. The admState dependence between UserObject and ResourceB must be checked at updating of the object ResourceB as well as at updating of UserObject instances. Therefore, this condition is implemented in the method checkConsistency in ResourceB.

Figure 9 shows the declaration file, which is generated from the specification of ResourceA. The opState dependence of ResourceB is implemented by derivation of the value of opState from resourceBstate and internalOpState, compare Figure 20.

The advantages of the above described are the following.

Dependences and co-operation between uncoordinated objects can be specified and implemented in the same way as dependences between objects in the same subsystem. This implies that they are visible to and can be interpreted by a managing system.

The above illustrated process for maintenance of dependences between uncoordinated objects in a controlled and visible way facilitates implementation of a determinable management model in a layered architecture.

The degree of reuseability of the library components can be improved by a high degree of flexibility when combinations of components are concerned.

Claims

1. A method for implementing a managed object in a subsystem of a managed system in a management network with at least one managing system and at least one managed system, for telecom or open systems, said managed system consisting of subsystems forming each a part of the managed system including one or more managed objects, characterized by implementing the managed objects in the subsystem, uncoordinatedly with respect to other subsystems, in such a way that they can be connected to and transmit messages to other objects in other subsystems, and without knowing the type of objects in the other subsystems, by designing a first object (330) for co-operation with an abstract object (332) defining an interface consisting of unimplemented methods, which may be called by the first object, and letting, at later design of a second object (334) unknown to said first object and intended to co-operate with the first object, the other object inherit the abstract object and implement the inherited methods, so that the first object at co-operation with the second object will consider it as being of said abstract type.

2. A method according to claim 1, characterized by delegating, at call of a method defined in the interface of the abstract object, the call to the implementation in the real object by means of late binding.

3. A method according to claims 1 or 2, characterized by reloading, in case of loading of said second object, said first object in the managed system by means of dynamic linking between the respective subsystems.

4. A method according to any of claims 1-3, characterized by locating said first and second objects in first and second subsystems, respectively.

5. A method according to claim 4, characterized by using dynamic linking to enable loading in the managed system of said second subsystem without reloading of said first subsystem.

Fig.1

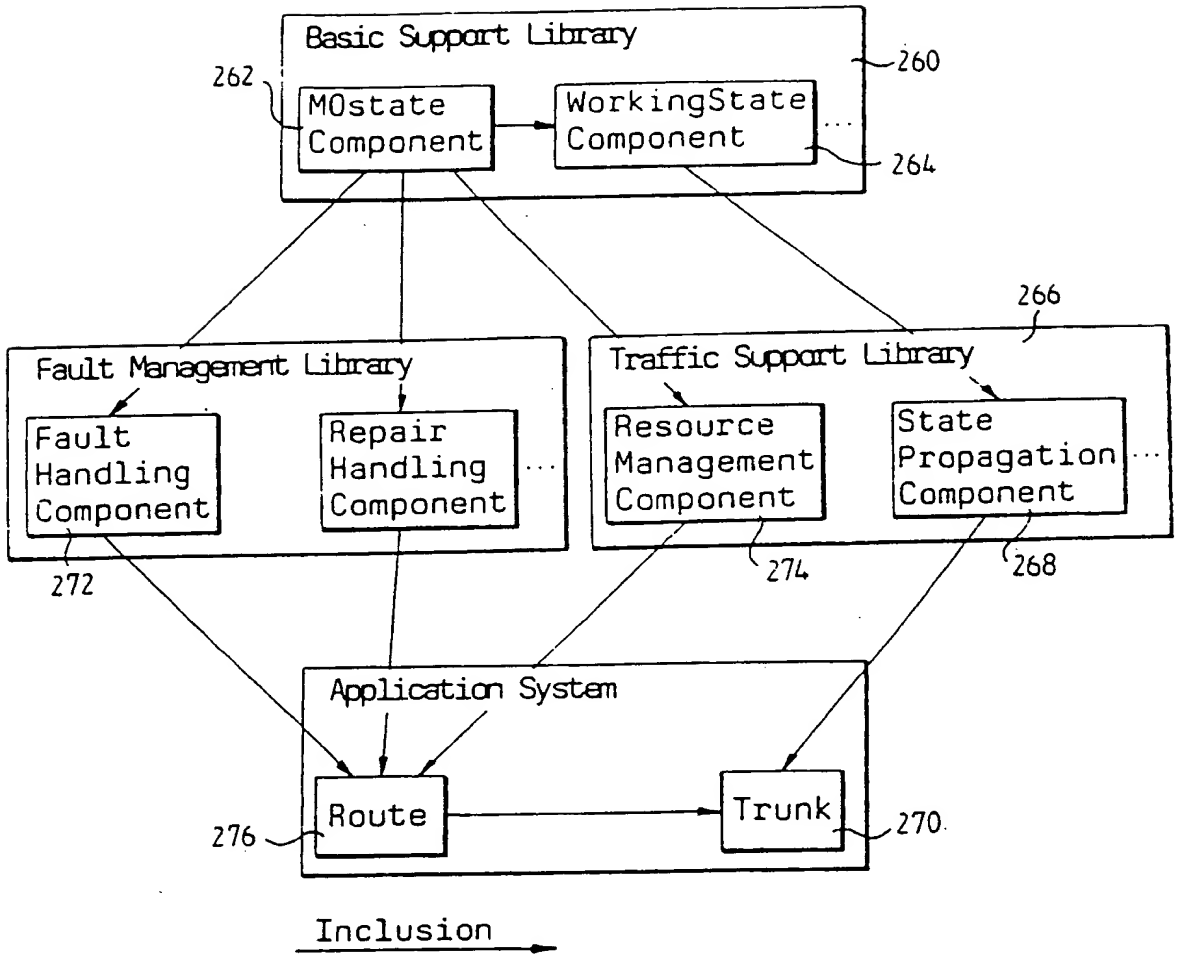

Basic Support Library
260
262
MOstate Component
WorkingState Component
264
266
Fault Management Library
Traffic Support Library
Fault Handling Component
Repair Handling Component
Resource Management Component
State Propagation Component
272
274
268
Application System
Route
Trunk
276
270
Inclusion

Fig.2

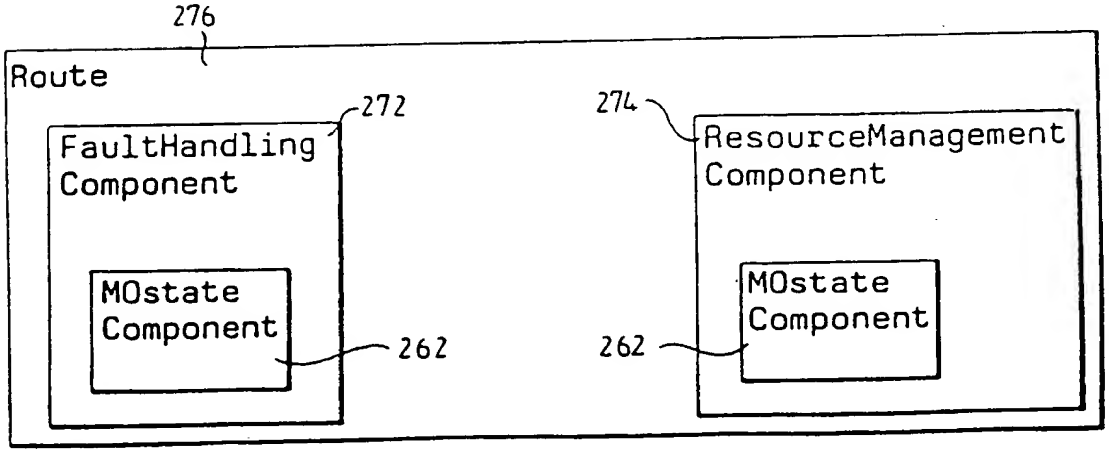

276
Route
272
FaultHandling Component
MOstate Component
262
274
ResourceManagement Component
MOstate Component
262

Fig.3

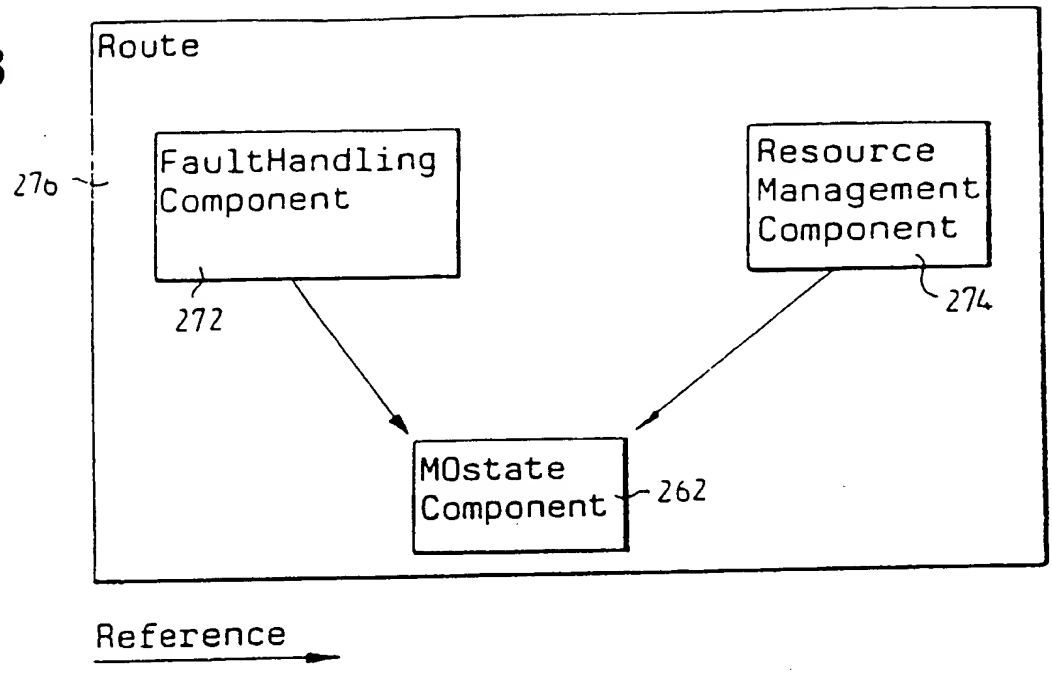


Fig.4

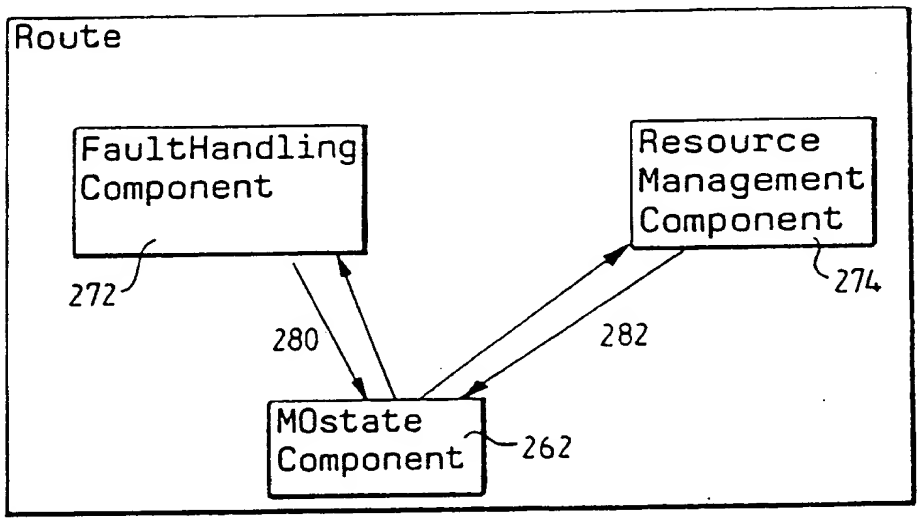


Fig.5

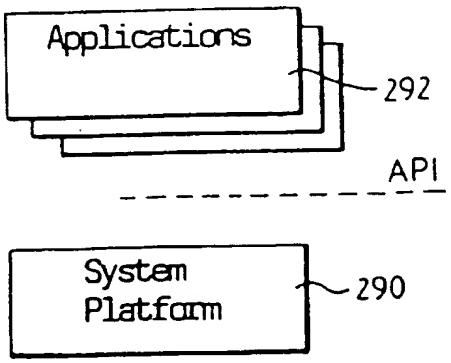


Fig.6

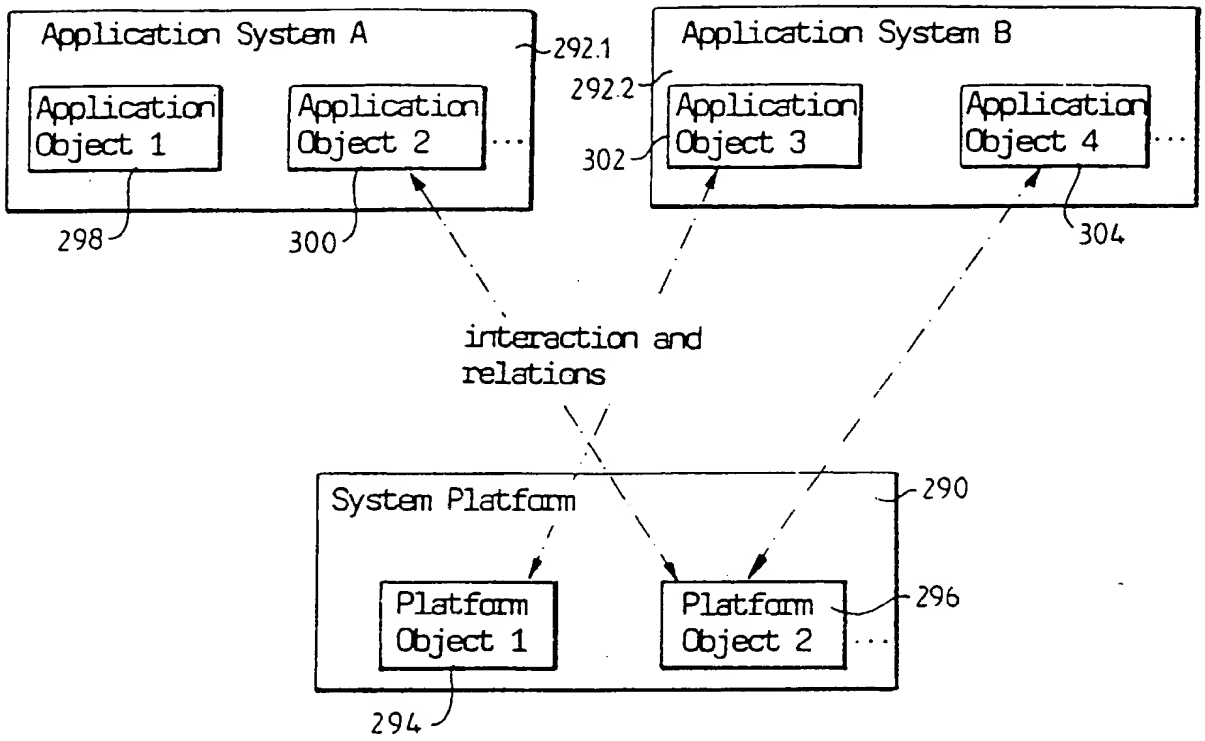

Application System A
292.1
Application Object 1
Application Object 2
298
300
Application System B
292.2
Application Object 3
Application Object 4
302
304
interaction and relations
System Platform
290
Platform Object 1
Platform Object 2
296
294

Fig.7

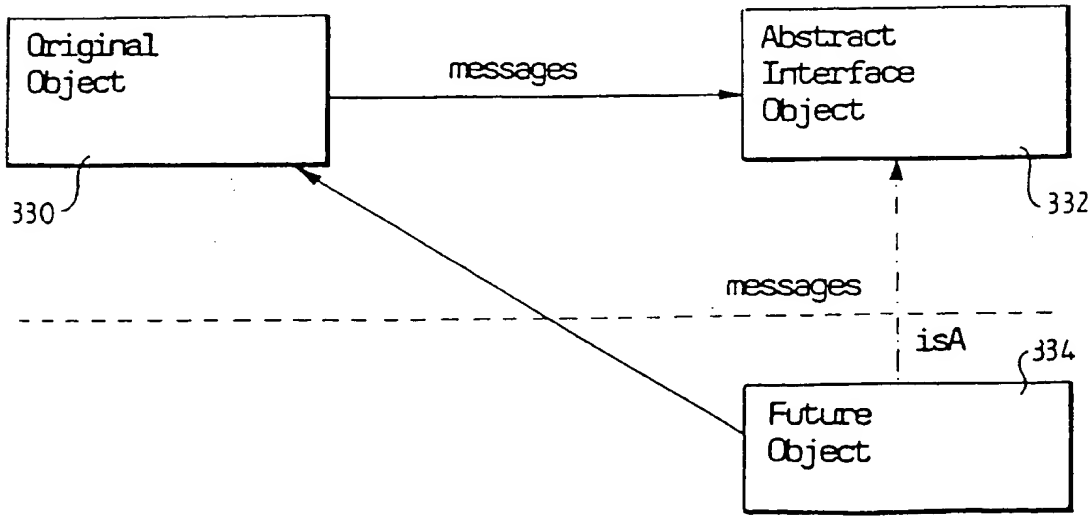

Original Object
messages
Abstract Interface Object
330
332
messages
isA
334
Future Object

Fig.8

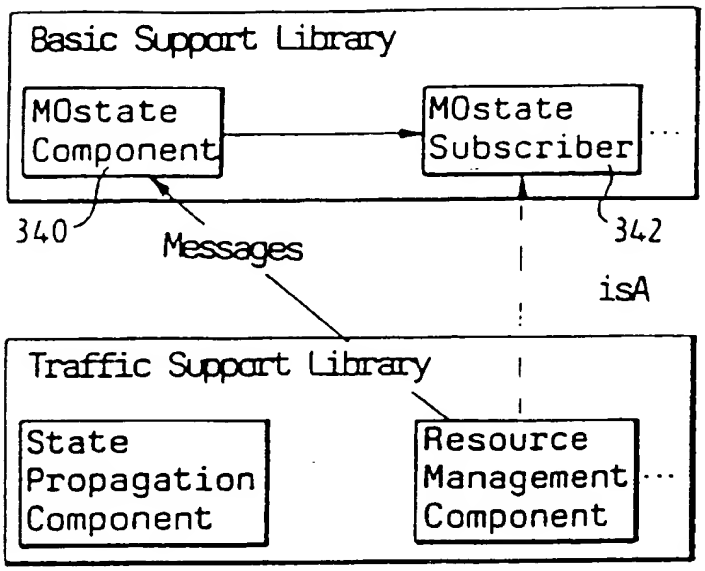


Fig.9

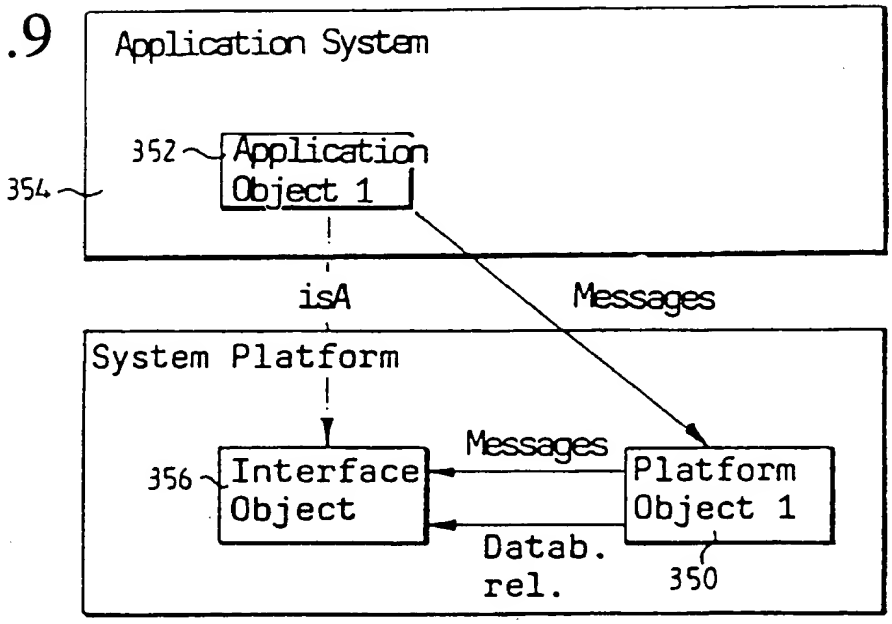


Fig.10

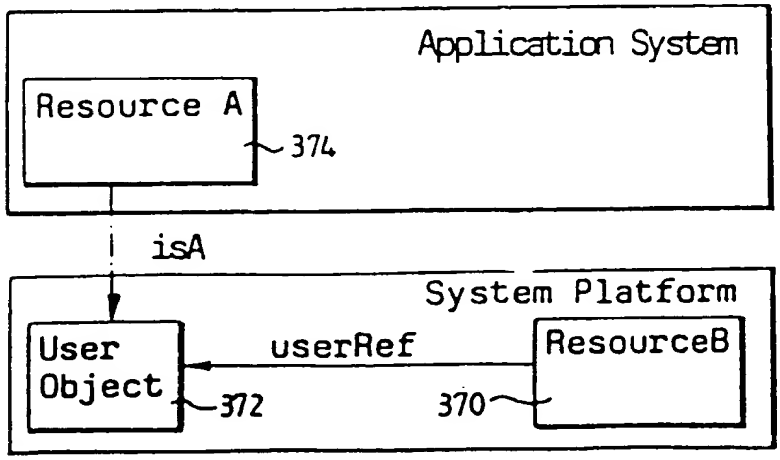


Fig.11

```
1  OBJECT TYPE ResourceB IS
2    ATTRIBUTES
3      key : KeyType;
4      admState : AdministrativeState;
5      opState : OperationalState;
6      userRef : REFERENCE TO UserObject
                 INVERSE OF Bref;
7    PRIMARY KEY key;
8
9    METHODS
10     allocate() RETURNS Boolean;
11
12   PRE-CONDITIONS
13     deleteObject() ONLY IF
       admState = locked;
14
15   PARTY TO ResourceAndUser;
16END;
```

Fig.12

```
17OBJECT TYPE UserObject IS
18   ATTRIBUTES
19      admState : AdministrativeState,
                   PURE VIRTUAL;
20      resourceBstate : OperationalState;
21      Bref : REFERENCE TO ResourceB
               INVERSE OF UserRef;
22
23   PARTY TO ResourceAndUser;
24END;
```

Fig.13

```
25DEPENDENCY SCHEMA ResourceAndUser IS
26   FOR ALL UserObject(a), ResourceB(b);
27   RELATIONS a.Bref = b;
28
29   POST-CONDITIONS
30     NOT (a.admState = unlocked AND
31           b.admState = locked)
32     RULES
33      WHEN COMMIT THEN CHECK CONSISTENCY;
34
35   PROPAGATIONS
36     WHEN b.opState = NewValue
37     THEN CONCLUDE a.resourceBstate=NewValue;
38
39END;
```

Fig.14

```
40OBJECT TYPE ResourceA IS
41   BASE UserObject;
42
43   ATTRIBUTES
44     key : KeyType;
45     admState : AdministrativeState;
46     DERIVED opState = IF (InternalOpState=disabled OR
47                           resourceBstate=disabled)
48                       THEN disabled
49                       ELSE enabled;
50
51     internalOpState : OperationalState PRIVATE;
52
53   PRIMARY KEY key;
54
55   METHODS
56     allocate() RETURNS Boolean;
57
58   PRE-CONDITIONS
59     deleteOblect() ONLY IF admState = locked;
60
61   POST-CONDITIONS
62     NOT (admState = unlocked AND Bref = NULL);
63
64END;
```

Fig.15

```
65 ifndef _UserObject_hh_
66 define _UserObject_hh_
67
68 include <PredefinedTypes.hh>
69 include "MOstateTypes.hh"
70
71class ResourceB
72
73// OBJECT Type Userobject
74
75class Userobject
76(
77public
78   static UserObject* open(Mode,const
79   KeyType&, DbTransaction* transaction=
     NULL);
80   virtual deleteObject() =0;
81   virtual Boolean checkConsistency
     (ErrorMessage&);
82
83   virtual AdministrativeState
     getAdmState() =0;
84   virtual void setAdmState
     (const AdministrativeState)==;
85   OperationalState getResourceBstate();
86   setResourceBstate(const
     OperationalState);
87   ResourceB* getBref();
88   void setBref(ResourceB*);
89
90private:
91.....
92.....
93);
94
95 endif
```

Fig.16

```
97 include "UserObject.hh"
98 include "ResourceB.hh"
99
100//////////////////////
101//
102// UserObject: METHOD checkConsistency
103//
104
105Boolean UserObject::checkConsistency(
            ErrorMessage& message)
106(
107   Boolean flag = true;
108
109   flag = !(getAdmState()  unlocked &&
110              getBref()->get admState()==locked);
111   if ( !flag) (
112      createErrorMessage(1,message);
113   )
114   return flag;
115)
116
```

Fig.17

```
117 ifndef __ResourceB_hh_
118 define _ResourceB_hh_
119
120 include <PredefinedTypes.hh>
121 include "MOstateTypes.hh"
122
123class UserObject;
124
125// OBJECT TYPE ResourceB
126
127class ResourceB
128(
129public
130   void deletobject();
131
132   void deleteObject();
133   Boolean allocate();
134   virtual Boolean checkConsistency
      (ErrorMessage&);
135
136   KeyType getKey();
137   AdministrativeState getAdmState();
138   void setAdmState (const
      AdministrativeState);
139   OperationalState getOpState();
140   void setOpState(const OperationalState);
141   UserObject* getUserRef();
142   void setUserRef(UserObject*);
143
144private
145   void propagateOpState
      (const OperationalState);
146   void opState (const OperationalState);
147   .....
148   .....
149);
150
151 endif
```

Fig.18

```
153 include "ResourceB.hh"
154 include "UserObject.hh"
155
156//////////////////////
157//
158// ResourceB: METHOD setOpState
159//
160
161void ResourceB: :setOpState(
            const OperationalState newValue)
162(
163   OperationalState oldValue = getOpState();
164
165   opState(newValue);
166   if (newValue != oldValue)
167   (
168      propagateOpState(newValue);
169   )
170)
171
172//////////////////////////
173//
174// ResourceB: METHOD PropagateOpState
175//
176
177void ResourceB: :propagateOpState(
              const OperationalState newValue)
178(
179   getUserRef()->resourceBstate(newValue)
180)
181
```

Fig.19

```
182 ifndef _ResourceA_hh_
183 define _ResourceA_hh_
184
185 include <PredefinedTypes.hh>
186 include "UserObject.hh"
187
188// OBJECT TYPE ResourceA
189
190class ResourceA : public UserObject
191(
192public
193   static ResourceA* open(Mode,
194   const KeyType&, DbTransaction*
      transaction=NULL);
195   void deleteObject ();
196   virtual Boolean checkConsistency
      (ErrorMessage&);
197
198   Boolean allocate();
199
200   KeyType getKey();
201   AdministrativeState getAdmState();
202   void setAdmState (const
      AdministrativeState);
203   OperationalState getOpState();
204
205private
206   OperationalState getInternalOpState();
207   void setInternalOpState (const
      OperationalState);
208   .....
209   .....
210);
211
212 endif
213
```

Fig.20

```
214 include "ResourceA.hh"
215
216//////////////////////
217//
218// ResourceA: METHOD getOpState
219//
220
2210perationalState ResourceA: :getOpState()
222(
223   if  (get InternalOpState() ==disabled
224         getResourceBstate() ==disabled)
225   (
226    return disabled;
227   )
228   else
229   (
230    return enabled;
231   )
232)
233
```

(19) Europäisches Patentamt
European Patent Office
Office européen des brevets

(11) EP 0 817 422 A3

(12) EUROPEAN PATENT APPLICATION

(88) Date of publication A3:
04.03.1998 Bulletin 1998/10

(43) Date of publication A2:
07.01.1998 Bulletin 1998/02

(21) Application number: **97117274.7**

(22) Date of filing: **18.08.1993**

(51) Int. Cl.[6]: **H04L 12/24**

(84) Designated Contracting States:
BE CH DE DK ES FR GB GR IE IT LI NL

(30) Priority: **28.08.1992 SE 9202488**
05.02.1993 SE 9300363

(62) Document number(s) of the earlier application(s) in accordance with Art. 76 EPC:
93919758.8 / 0 627 143

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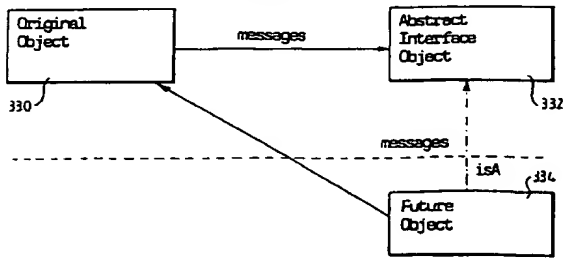
(54) A method for implementing a managed object in a network management system

(57) The invention relates to a method for implementing a managed object in a subsystem of a managed system in a management network with at least one managing system and at least one managed system, for telecom or open systems, said managed system consisting of subsystems forming each a part of the managed system including one or more managed objects.

The managed objects are implemented in the subsystem, uncoordinatedly with respect to other subsystems, in such a way that they can be connected to and transmit messages to other objects in other subsystems, and without knowing the type of objects in the other subsystems. This is obtained by designing a first object (330) for co-operation with an abstract object (332) defining an interface consisting of unimplemented methods, which may be called by the first object, and letting, at later design of a second object (334) unknown to said first object and intended to co-operate with the first object, the other object inherit the abstract object and implement the inherited methods, so that the first object at co-operation with the second object will consider it as being of said abstract type.

Fig.7



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DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int. Cl.6)
A	US 4782506 A (MAXIMILIAN SEVCIK), 1 November 1988 (01.11.88) * column 1, line 51 - column 3, line 53, figures 1,2 *	1-5	H04L 12/24
A	EP 0442809 A2 (DIGITAL EQUIPMENT CORPORATION), 21 August 1991 (21.08.91) * page 2, line 21 - line 29; page 3, line 28 - line 55, figure 1 *	1-5	
A	US 4903263 A (RAJENDRA PATEL ET AL), 20 February 1990 (20.02.90) * column 1, line 8 - line 21; column 12, line 49 - line 65, figure 2 *	1-5	

TECHNICAL FIELDS SEARCHED (Int. Cl.6)

H04L
H04M
H04Q

The present search report has been drawn up for all claims

Place of search	Date of completion of the search	Examiner
STOCKHOLM	14 November 1997	ANDERS STRÖBECK

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